

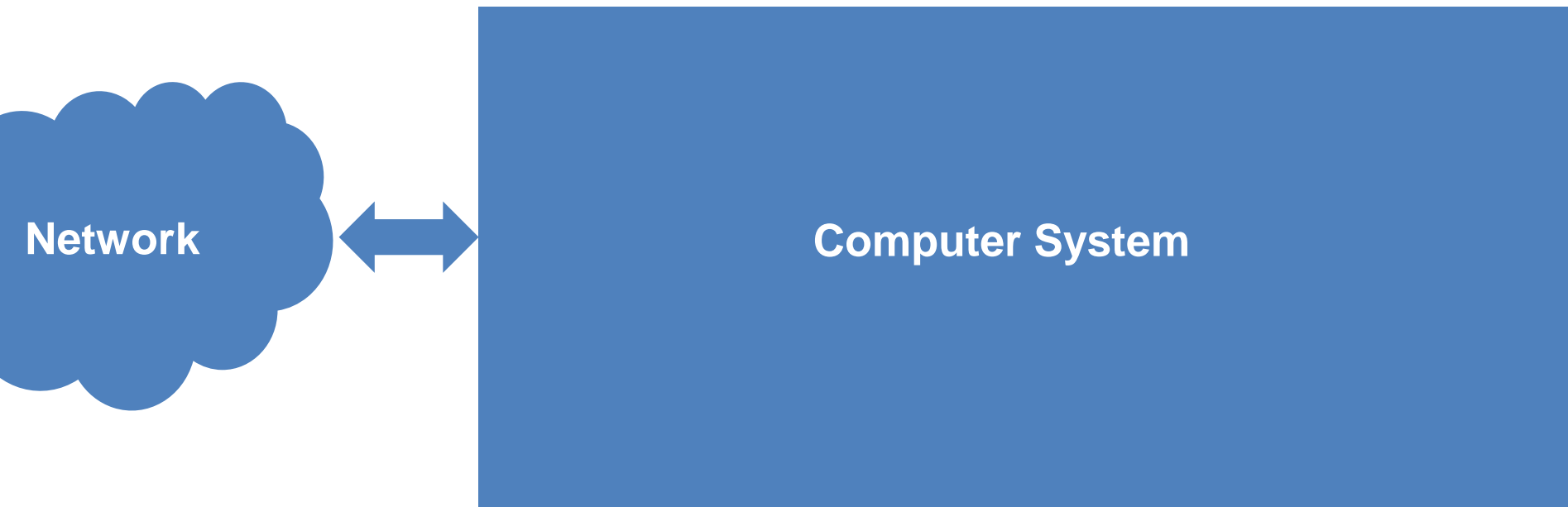
Protecting RISC-V Processors against Physical Attacks

Stefan Mangard, Robert Schilling, Thomas Unterluggauer, Mario Werner

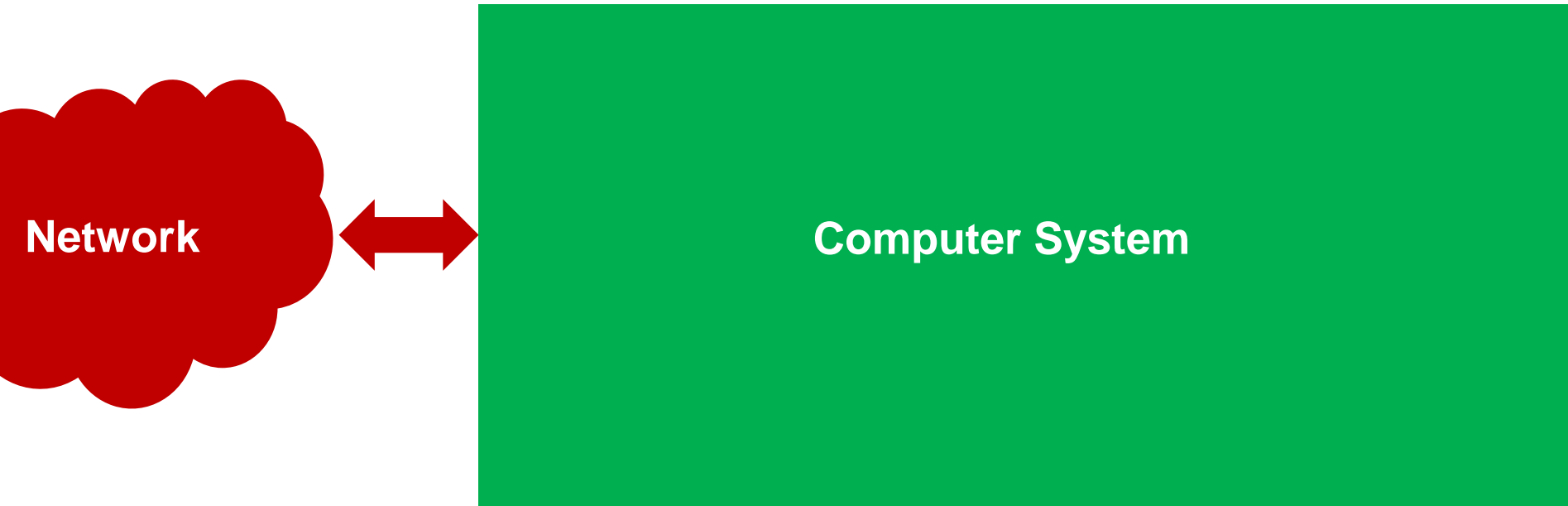
Graz University of Technology

March 28th, 2019

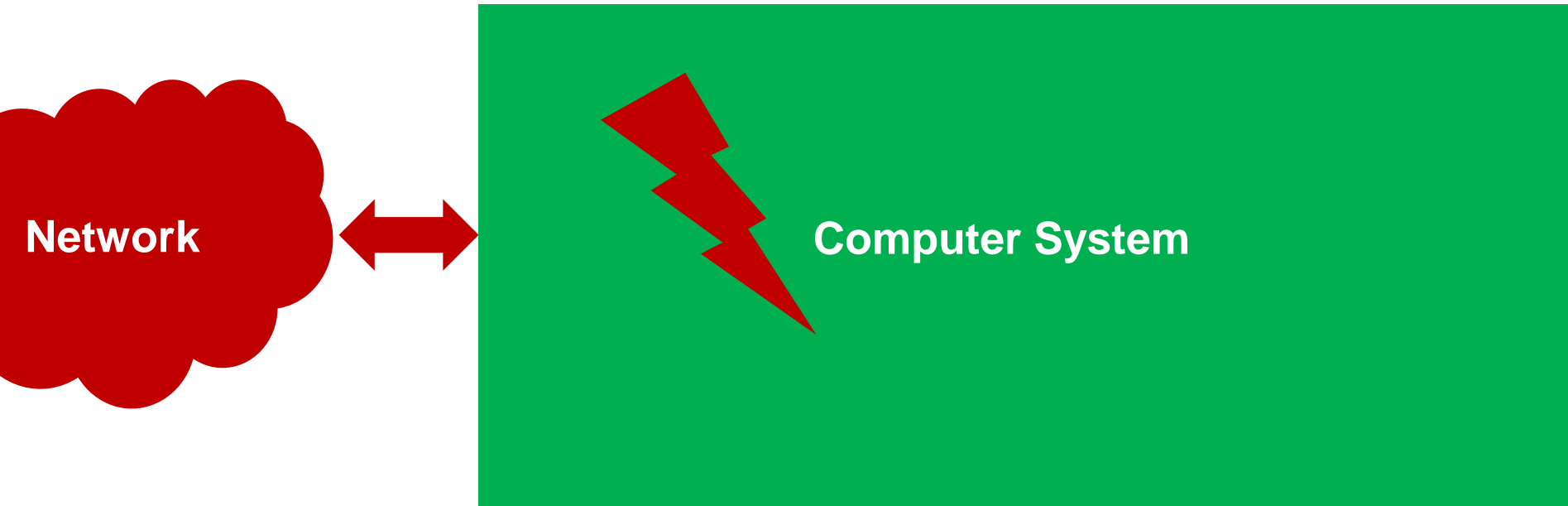
Attack Settings



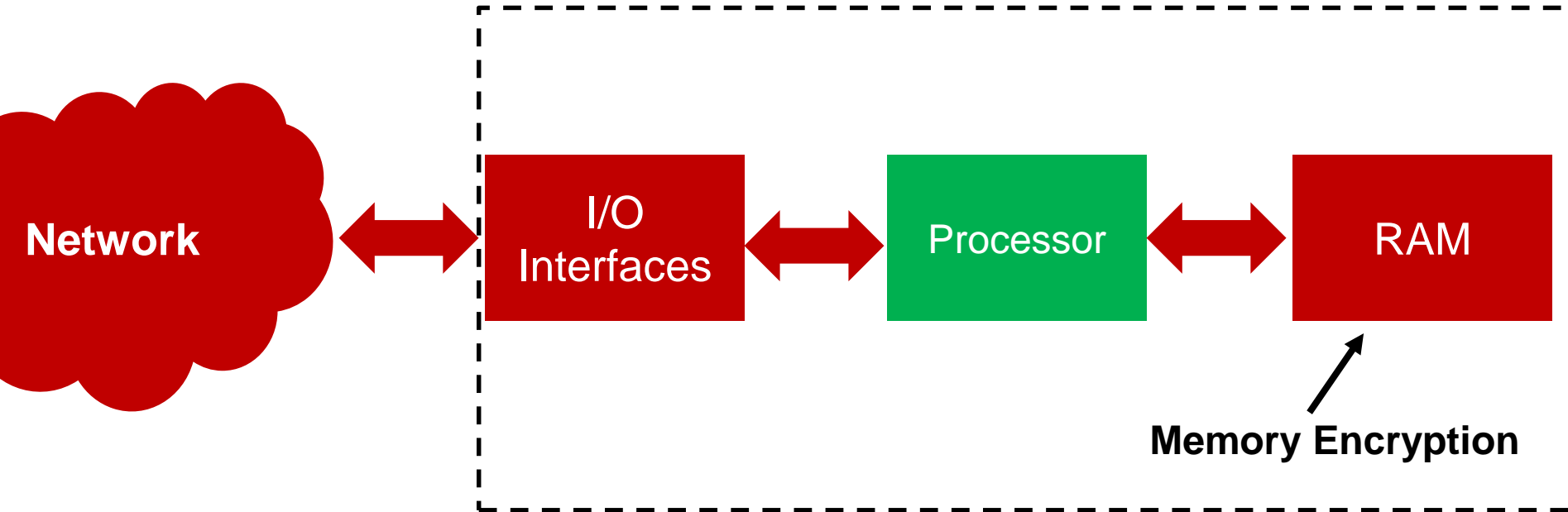
Classic Setting: Attacks Via The (Network) Interfaces



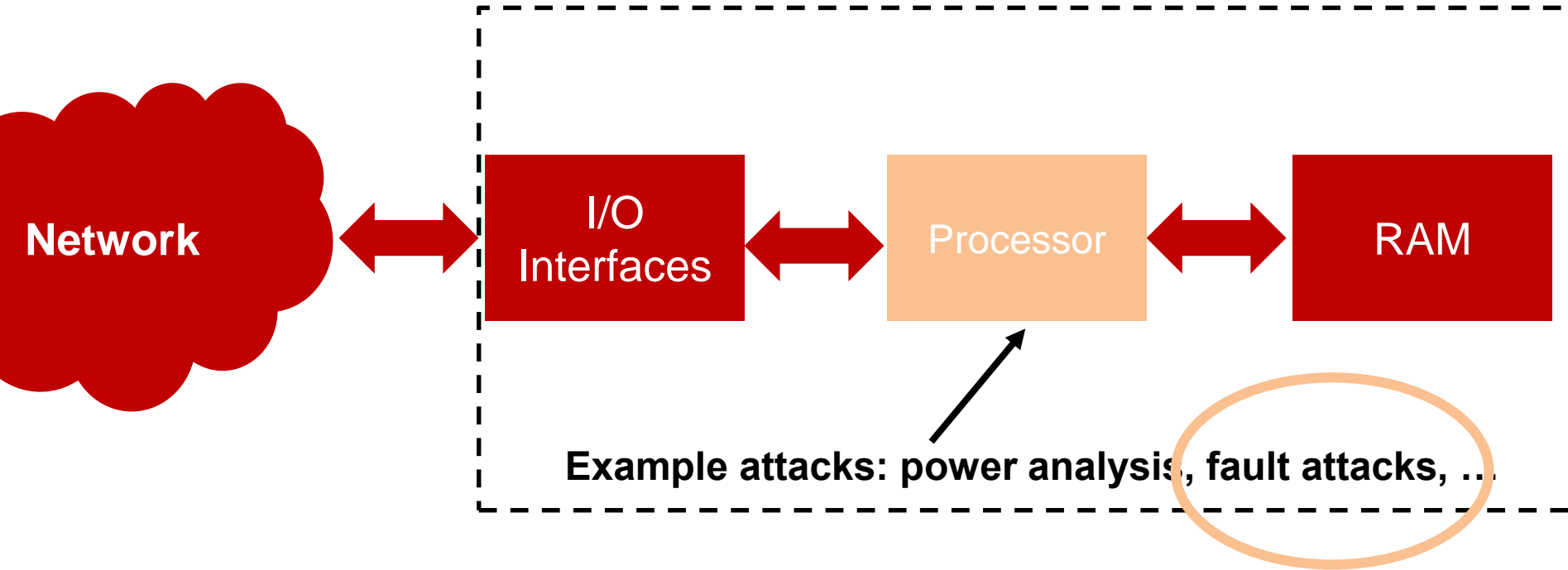
Classic Setting: Attacks Via The (Network) Interfaces



The Secure Processor Setting



Physical Attack Setting: Also Processor is Attacked



Small Faults – Big Effects

1 bit

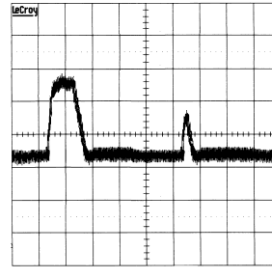
- determines your privilege level
- determines your access rights
- determines whether a password is considered correct or not
- Changes completely the meaning of an opcode
- ...

Fault Attacks - Relevance

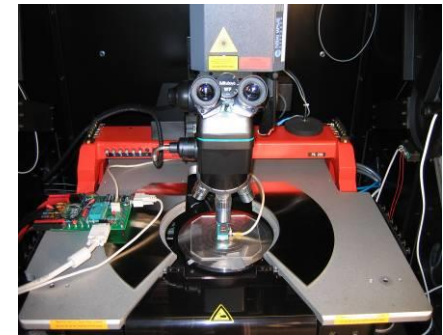
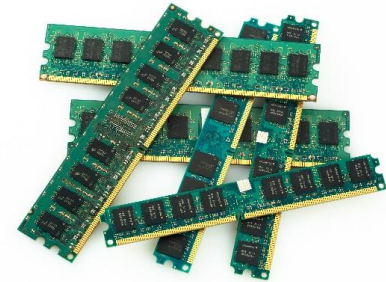
- Traditional fields
 - secure elements
 - smart cards, ...
- New fields:
 - IoT, automotive, etc.
 - Example: Privilege escalation on Linux (Niek Timmers, Cristofaro Mune: Escalating Privileges in Linux Using Voltage Fault Injection. FDTC 2017: 1-8)

Fault Attacks – How?

- Most popular techniques
 - Voltage glitches
 - Laser

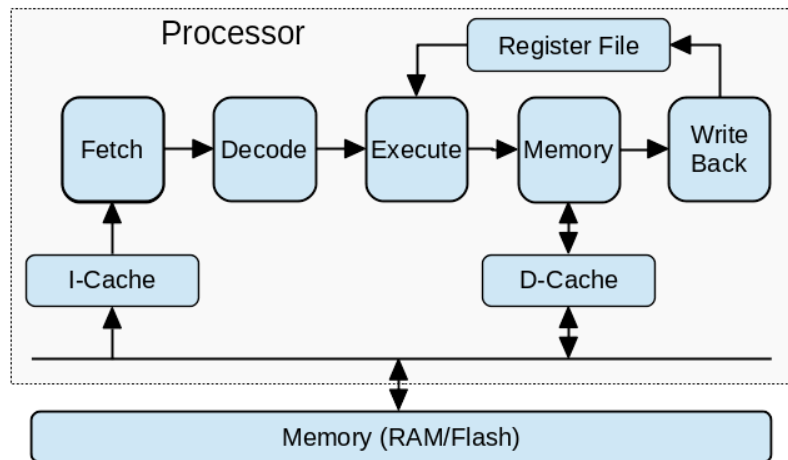


Rowhammer



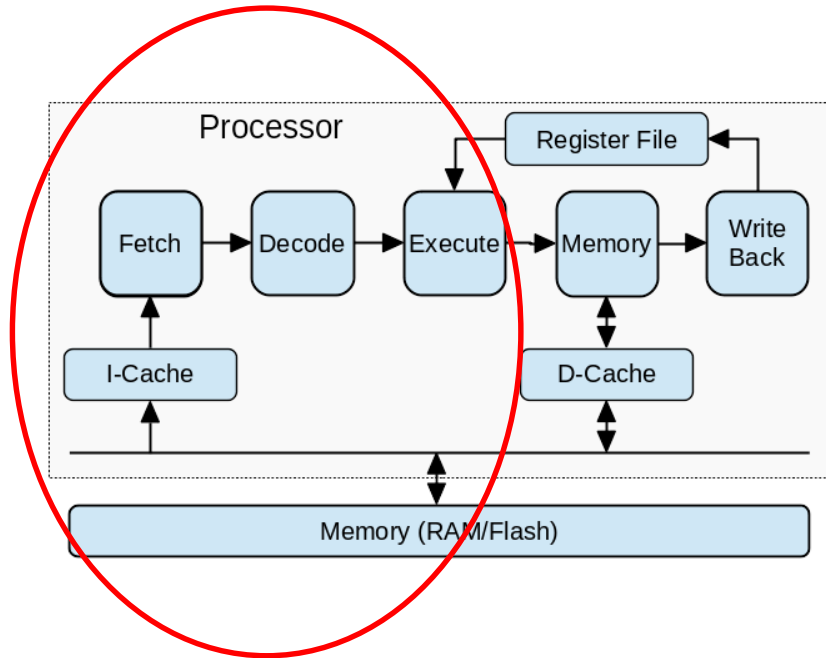
- Effects
 - Changes in data, pointers, ...
 - Changes in program flow: skip of instruction pointer, induction of branches, ...

Fault Attacks – What Can We Do?



Redundancy

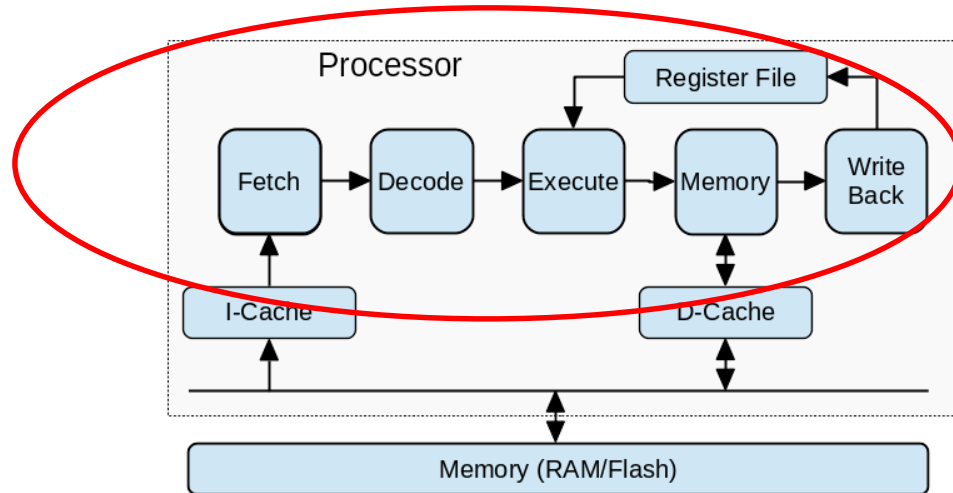
Fault Attacks – What Can We Do?



- **Control flow graph integrity**

Mario Werner, Thomas Unterluggauer, David Schaffenrath, Stefan Mangard. “Sponge-Based Control-Flow Protection for IoT Devices”. In: Euro S&P 2018

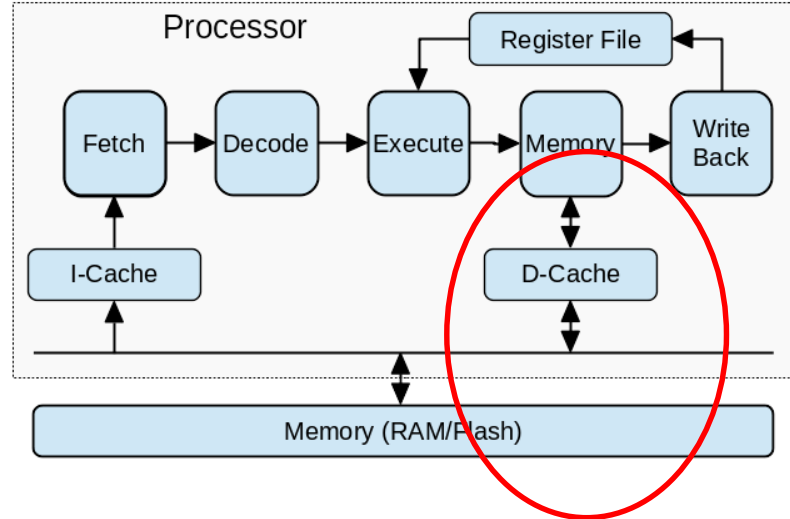
Fault Attacks – What Can We Do?



- **Branch Integrity**

Robert Schilling, Mario Werner, Stefan Mangard. “Securing Conditional Branches in the Presence of Fault Attacks”. In: DATE 2018

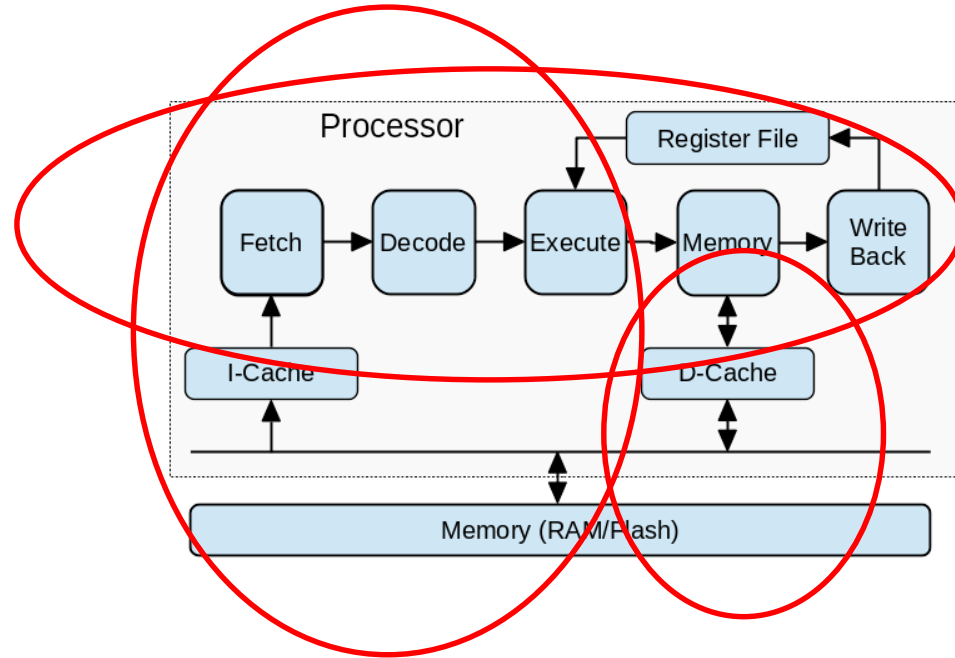
Fault Attacks – What Can We Do?



- **Memory Access**

Robert Schilling, Mario Werner, Pascal Nasahl, Stefan Mangard. “Pointing in the Right Direction-Securing Memory Accesses in a Faulty World”. In: ACSAC 2018

Fault Attacks – What can we do?

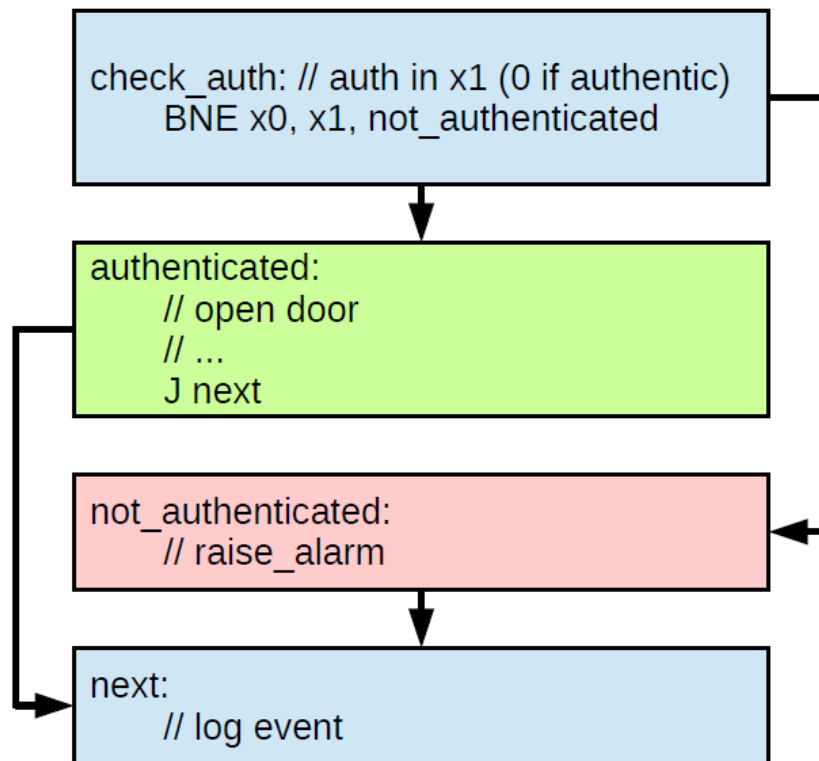


Sponge-Based Control-Flow Protection

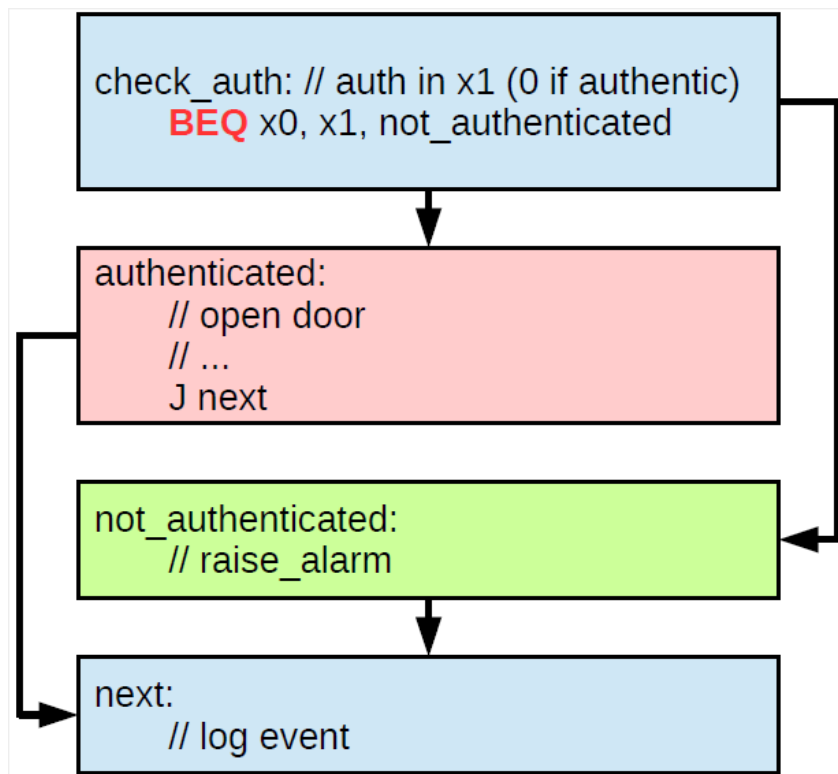
Control Flow Integrity is Critical

```
unsigned pin = read_pin();
bool auth = check_pin(pin);
if( auth ) {
    open_door();
} else {
    raise_alarm();
}
log_event();
```

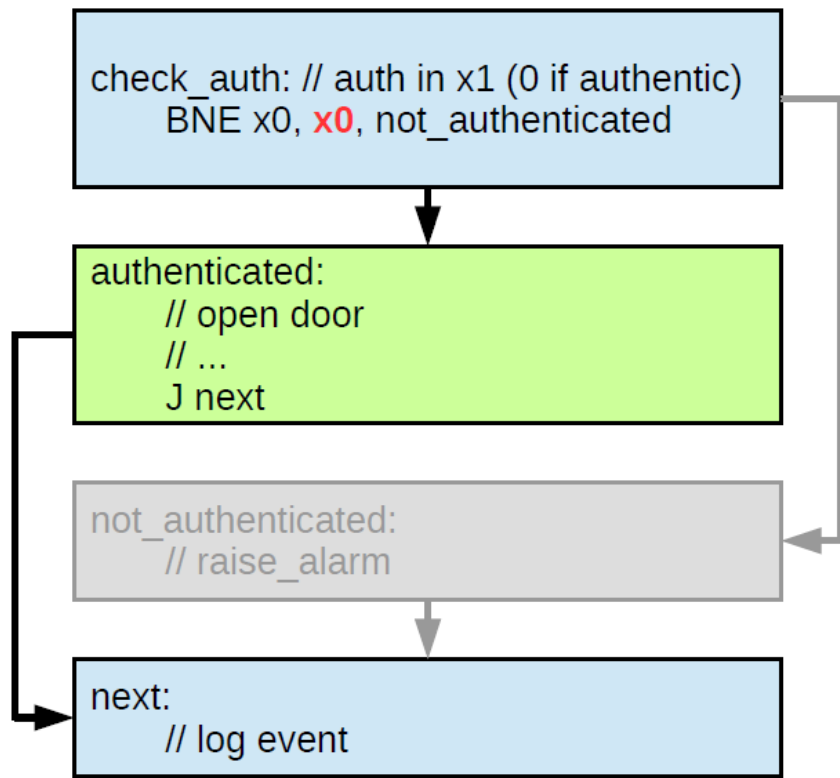

Control Flow Integrity is Critical



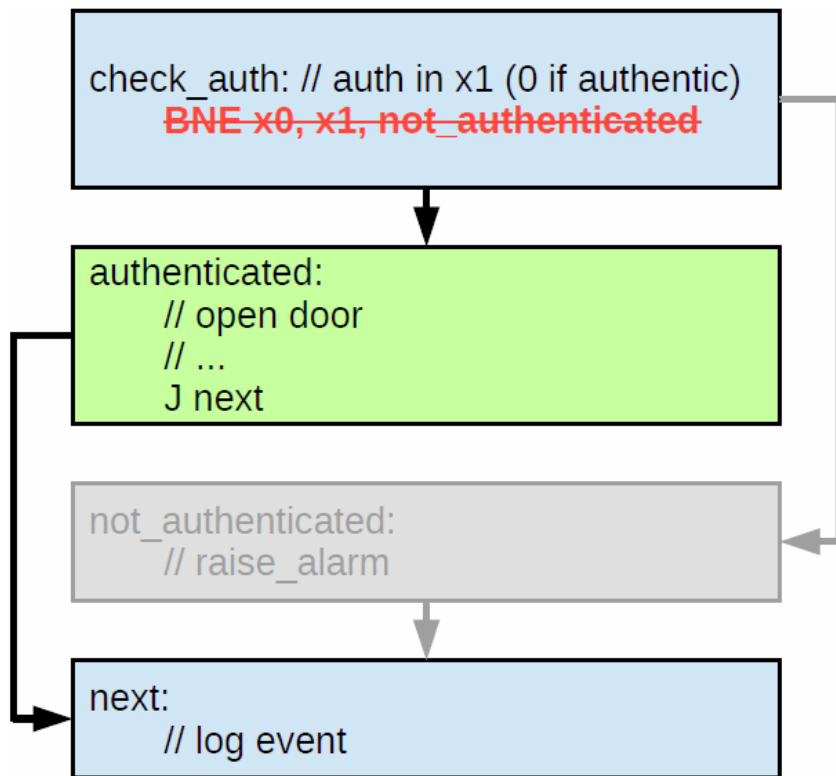
Control Flow Integrity is Critical



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Control Flow Integrity is Critical

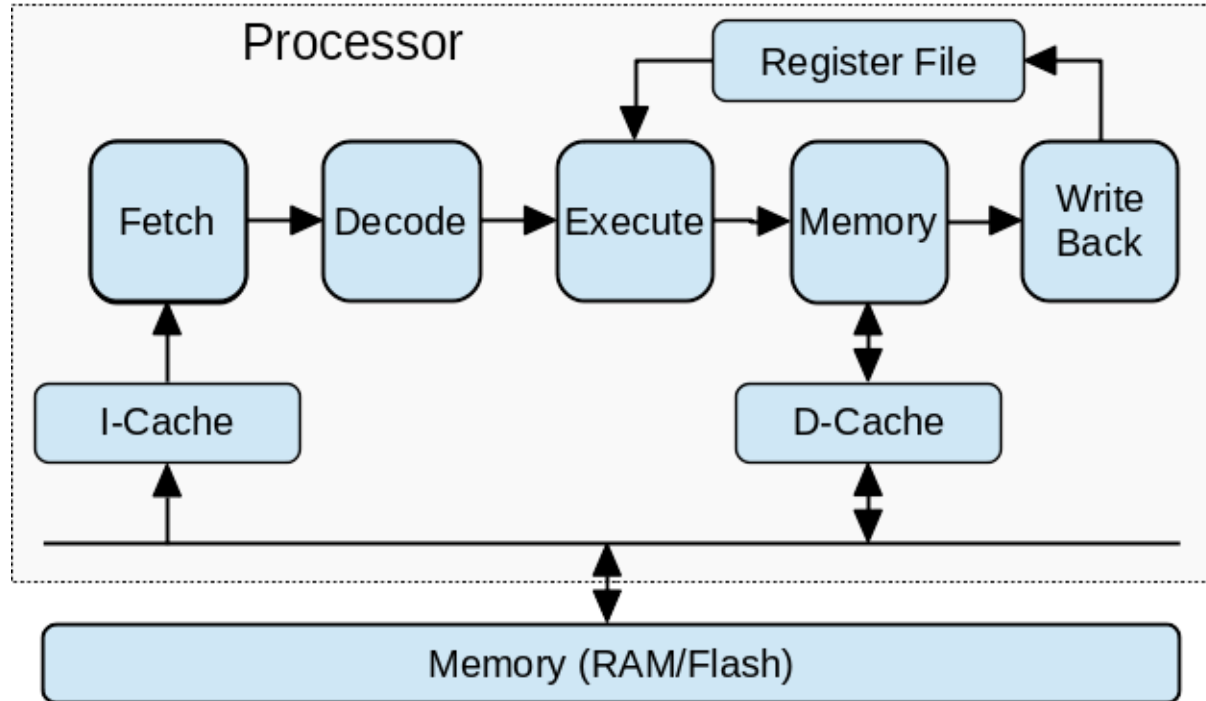


Main Idea

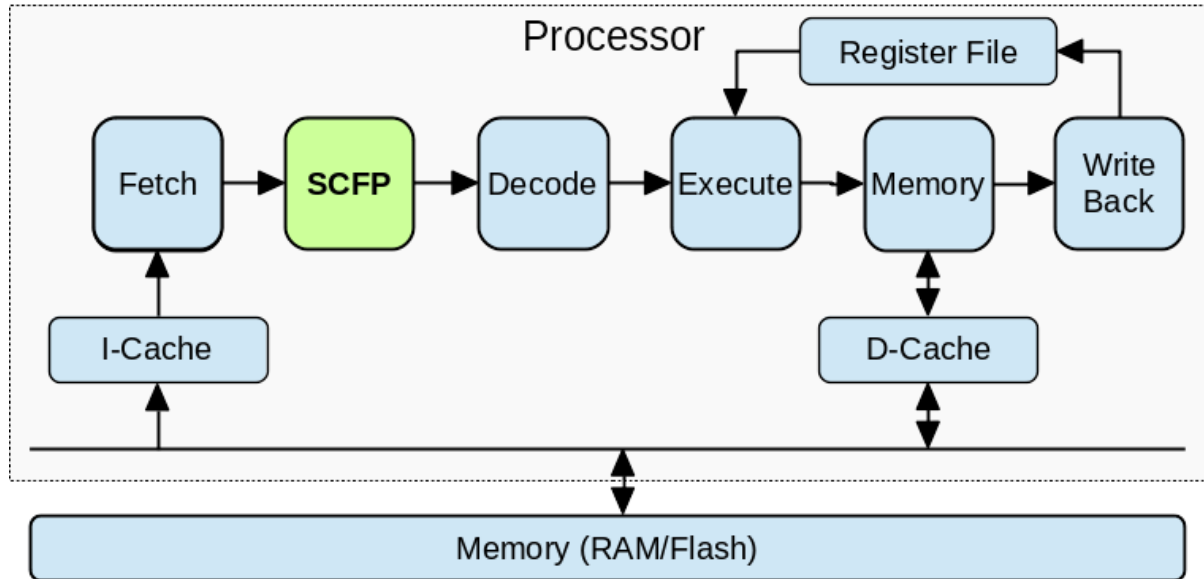
- Sponge-based Control-Flow Protection (SCFP)
 - Hardware supported CFI scheme
 - Encrypts the instruction stream with small granularity
 - Program can only be decrypted correctly as long it is executed correctly

- Costs
 - Highly configurable in terms of security and cost
 - RV32IM AEE-Light: ~10% runtime, ~20% code size

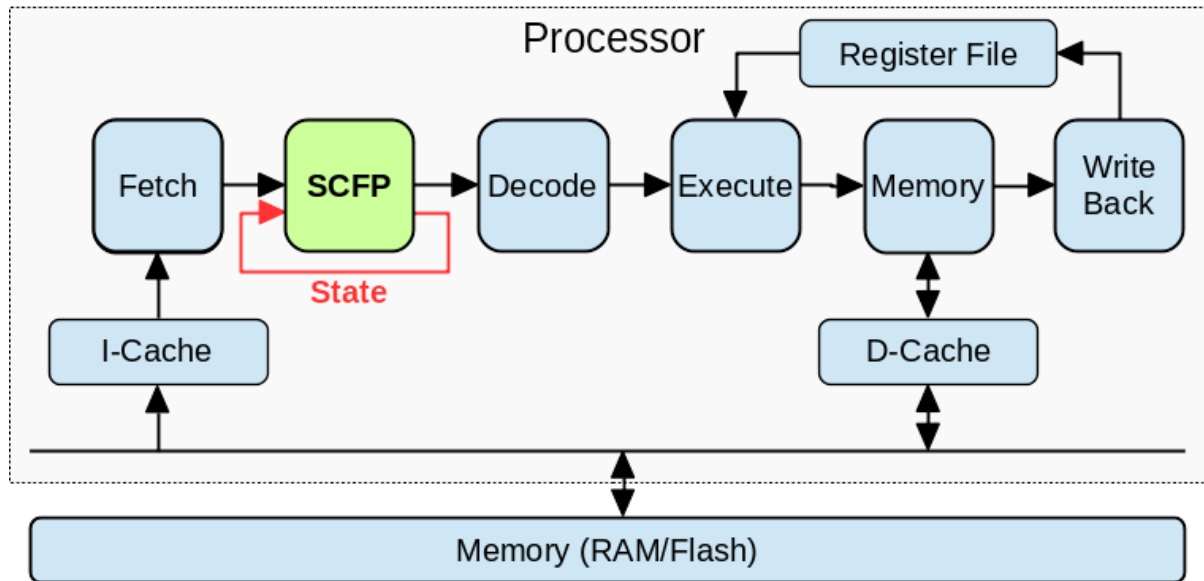
High Level Concept



High Level Concept



High Level Concept



Decryption/Execution Example

```
strcmp
```

```
: ec d0 ee 97  
: 28 ce 77 80  
: 75 41 64 b1
```

```
: 4b f4 51 75  
: d9 a6 02 ad  
: 51 7d 34 43
```

```
: 4d 1b c0 0f  
: a3 0f 21 3e
```

Decryption/Execution Example

```
strcmp  
0x1b2a0645  
: ec d0 ee 97  
: 28 ce 77 80  
: 75 41 64 b1
```

```
: 4b f4 51 75  
: d9 a6 02 ad  
: 51 7d 34 43
```

```
: 4d 1b c0 0f  
: a3 0f 21 3e
```

Decryption/Execution Example

```
strcmp  
0x1b2a0645  
0xdd3fbcce : 03 06 05 00 : 1b a2, 0(a0)  
              : 28 ce 77 80  
              : 75 41 64 b1
```

```
              : 4b f4 51 75  
              : d9 a6 02 ad  
              : 51 7d 34 43
```

```
              : 4d 1b c0 0f  
              : a3 0f 21 3e
```

Decryption/Execution Example

```

strcmp
0x1b2a0645
0xdd3fbcce : 03 06 05 00 : 1b a2, 0 (a0)
0xf5a92604 : 83 86 05 00 : 1b a3, 0 (a1)
           : 75 41 64 b1

```

```

           : 4b f4 51 75
           : d9 a6 02 ad
           : 51 7d 34 43

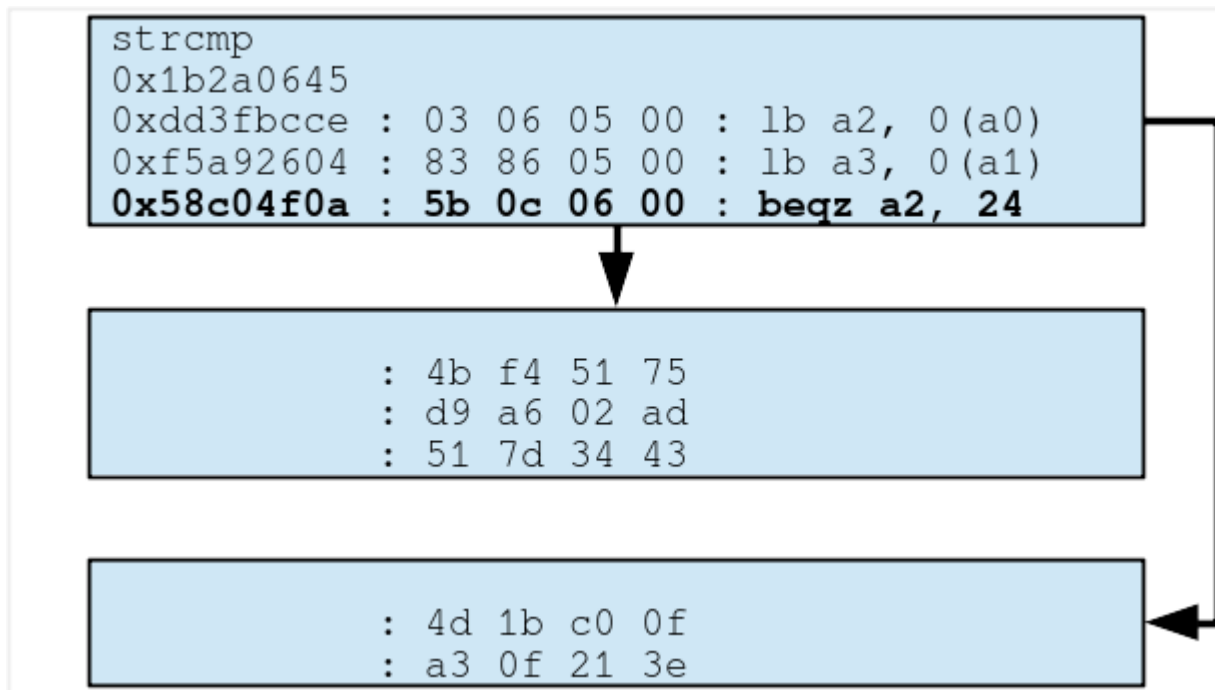
```

```

           : 4d 1b c0 0f
           : a3 0f 21 3e

```

Decryption/Execution Example



Decryption/Execution Example

```

strcmp
0x1b2a0645
0xdd3fbcce : 03 06 05 00 : lb a2, 0(a0)
0xf5a92604 : 83 86 05 00 : lb a3, 0(a1)
0x58c04f0a : 5b 0c 06 00 : beqz a2, 24

```

↓

```

0x58c04f0a
: 4b f4 51 75
: d9 a6 02 ad
: 51 7d 34 43

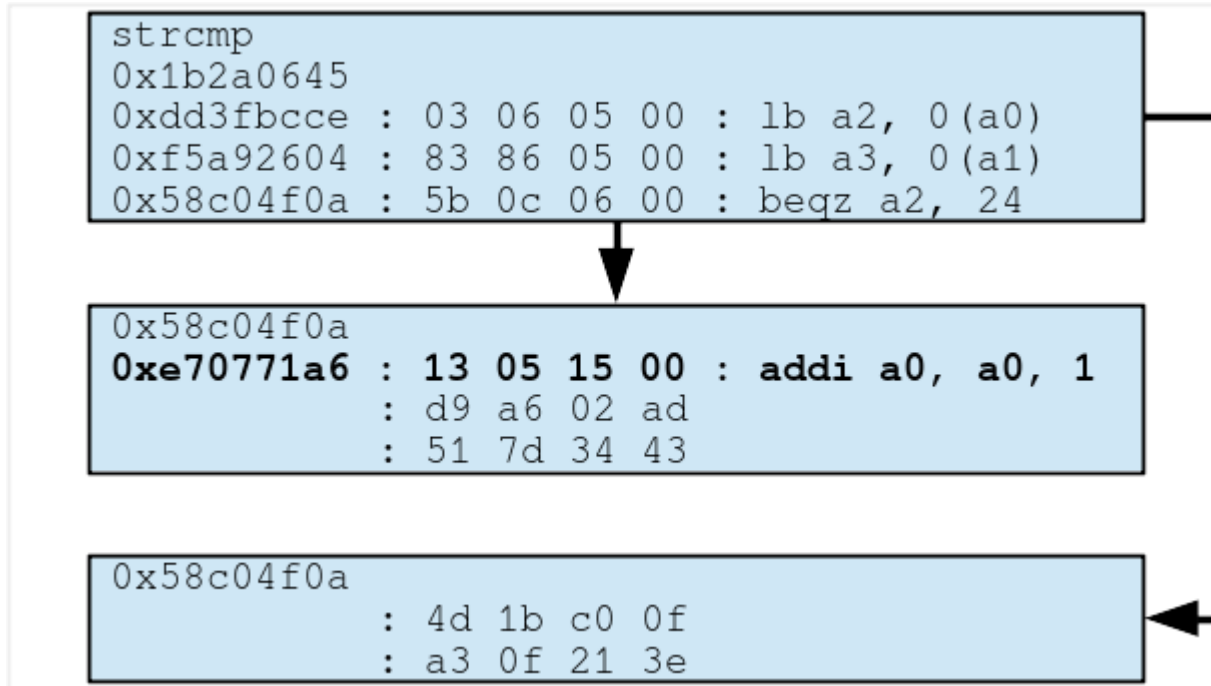
```

```

0x58c04f0a
: 4d 1b c0 0f
: a3 0f 21 3e

```

Decryption/Execution Example



Decryption/Execution Example

```

strcmp
0x1b2a0645
0xdd3fbcce : 03 06 05 00 : lb a2, 0(a0)
0xf5a92604 : 83 86 05 00 : lb a3, 0(a1)
0x58c04f0a : 5b 0c 06 00 : beqz a2, 24

```

```

0x58c04f0a
0xe70771a6 : 13 05 15 00 : addi a0, a0, 1
0x5b26165e : 93 85 15 00 : addi a1, a1, 1
           : 51 7d 34 43

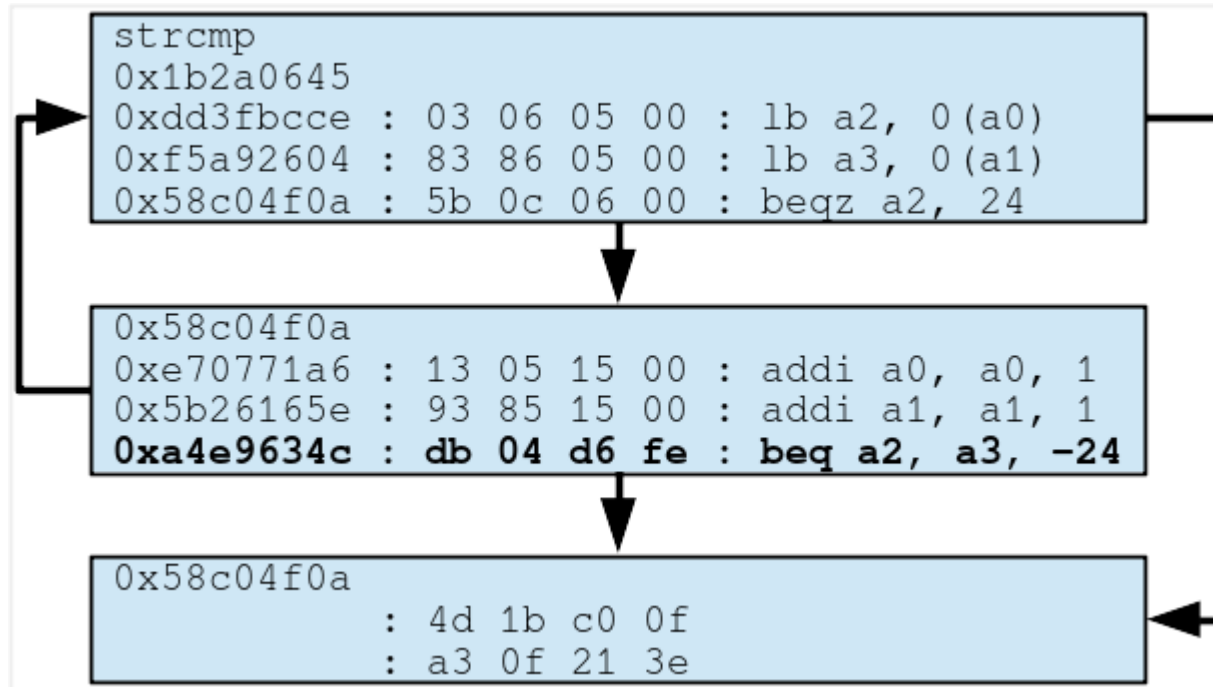
```

```

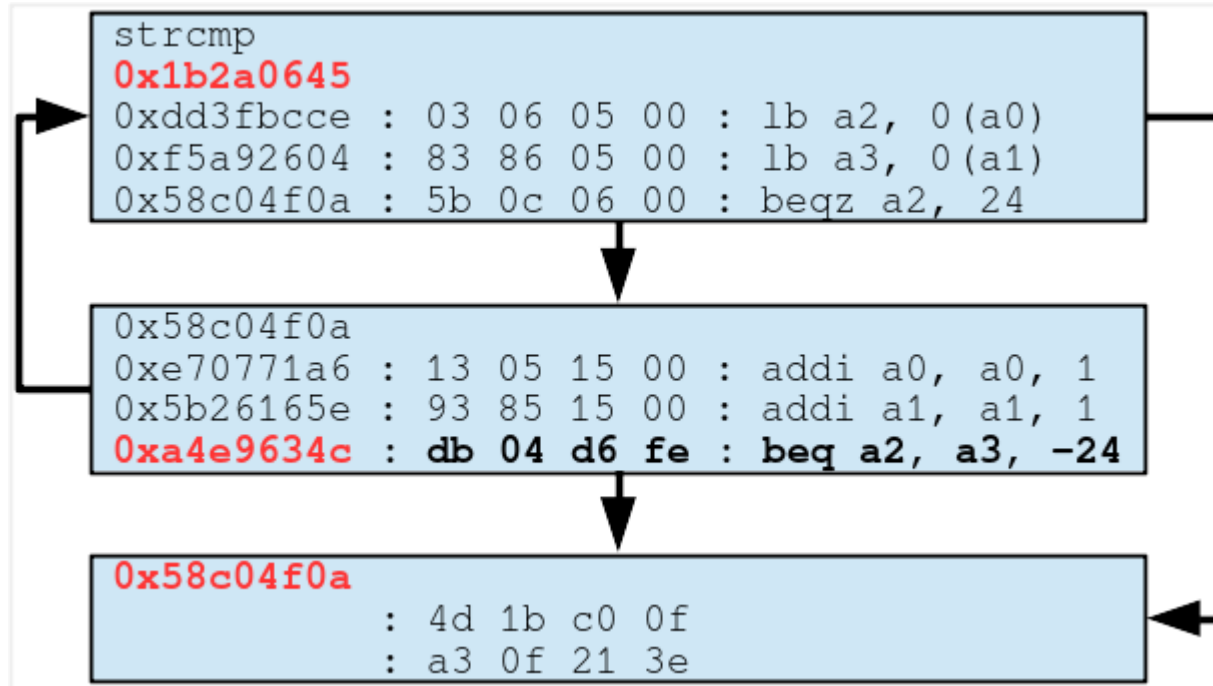
0x58c04f0a
           : 4d 1b c0 0f
           : a3 0f 21 3e

```

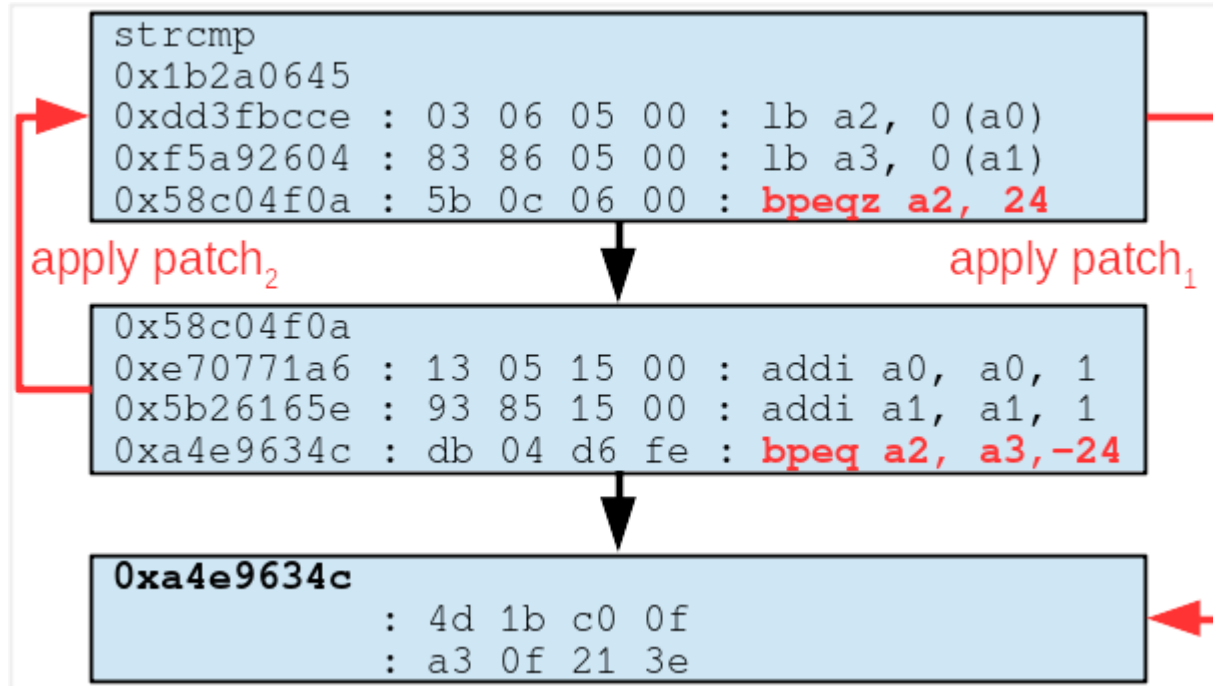

Decryption/Execution Example



Decryption/Execution Example

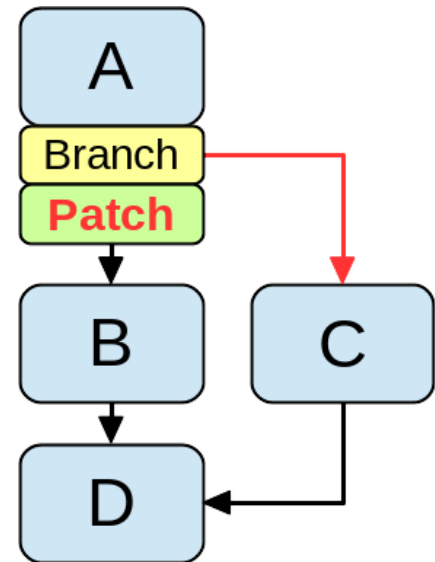


Decryption/Execution Example



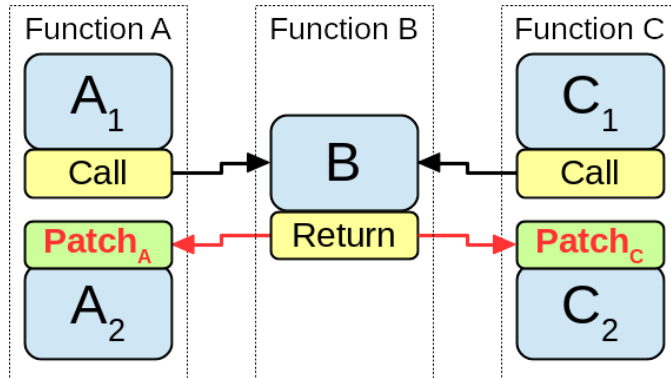
RISC-V ISA Integration - Branches

- Branches additional have an associated patch that is applied conditionally
- New BPEQ, BPNE, BPLT, BPLTU, BPGE, and BPGEU instructions

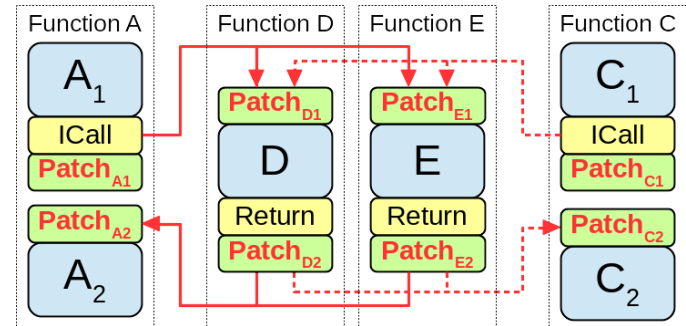


RISC-V ISA Integration – Calls

Direct Calls

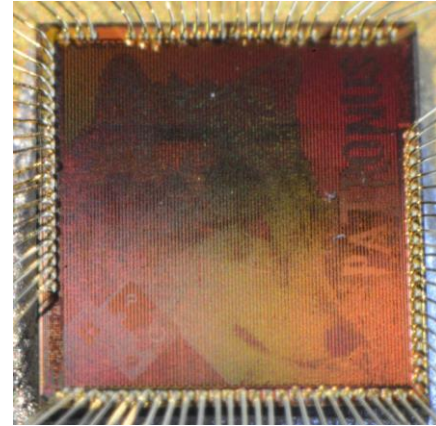


Indirect Calls



Prototype Implementation

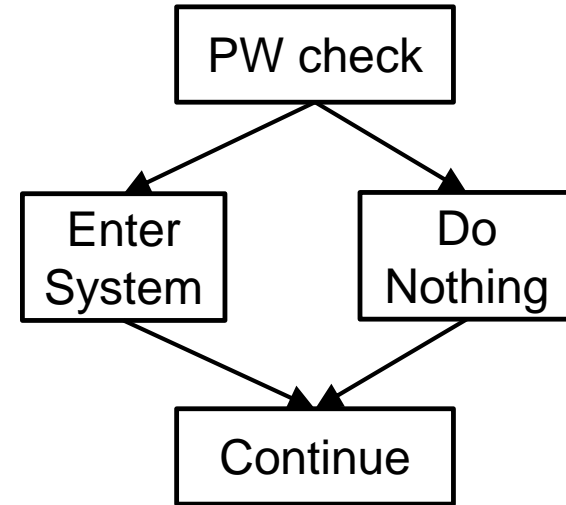
- LLVM-based toolchain
- RI5CY-based hardware
 - AEE-Light with PRINCE in APE-like mode
 - ~30kGE of area for SCFP at 100MHz in UMC65
 - ~10% runtime and ~20% code size overhead



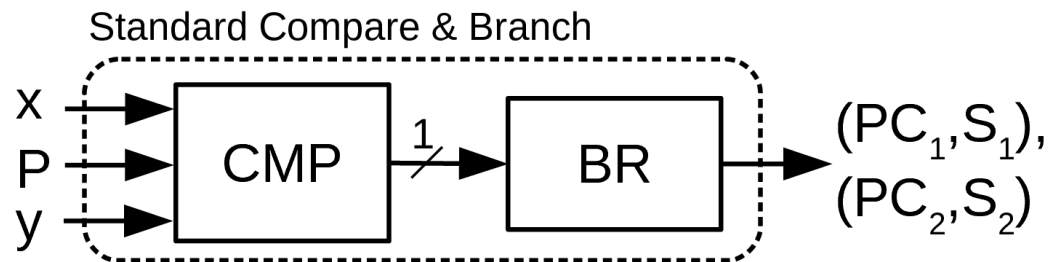
Protected Conditional Branches

Motivation

- Control-flow integrity (CFI) measures restrict the control-flow to valid execution traces
- Branching decisions require extra protection
- Examples of critical applications
 - Password checks, signature verification, ...



Classical Conditional Branch

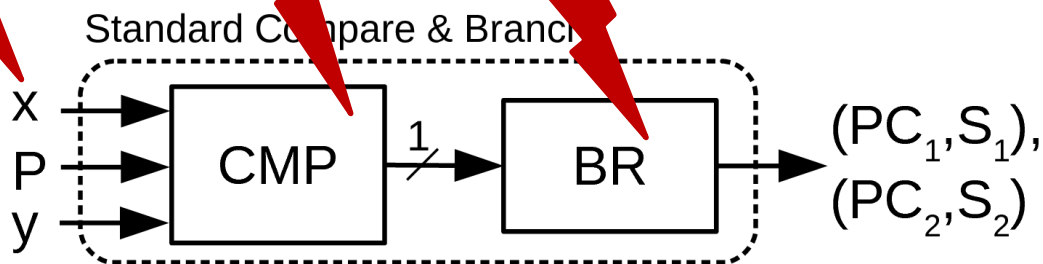


Faulting Conditional Branches

Fault the comparison

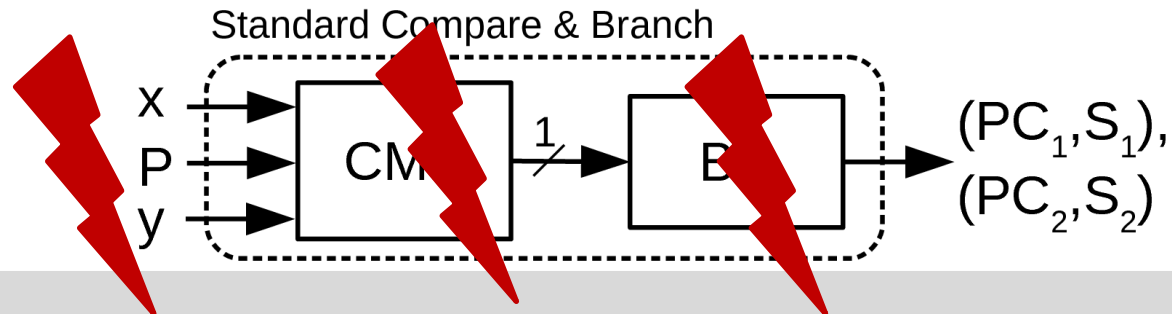
Fault the operands

Faulting the branch



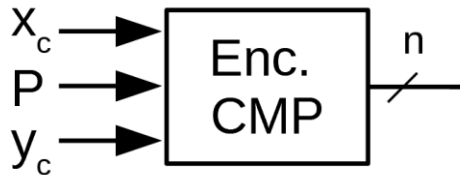
Design Goals

- Maximum flexibility concerning redundancy encoding of data (x, y)
 - Arithmetic codes are efficient for number encoding
 - Linear codes are used for strings
- Minimal changes to hardware



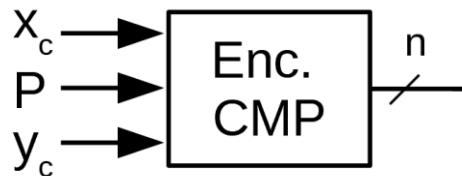
Step 1: Encoded Comparison in Software

- Efficiently possible e.g. for AN Codes
- Output of comparison: n-bit symbol for true or n-bit symbol for false



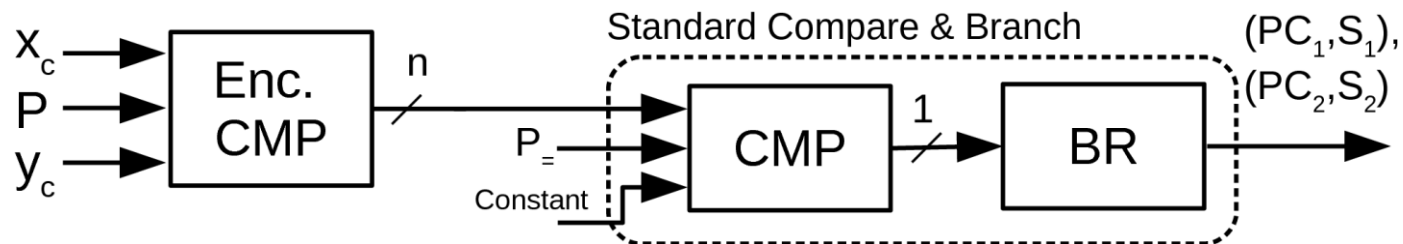
Step 1: Encoded Comparison in Software

How to securely branch based on the n-bit symbol?



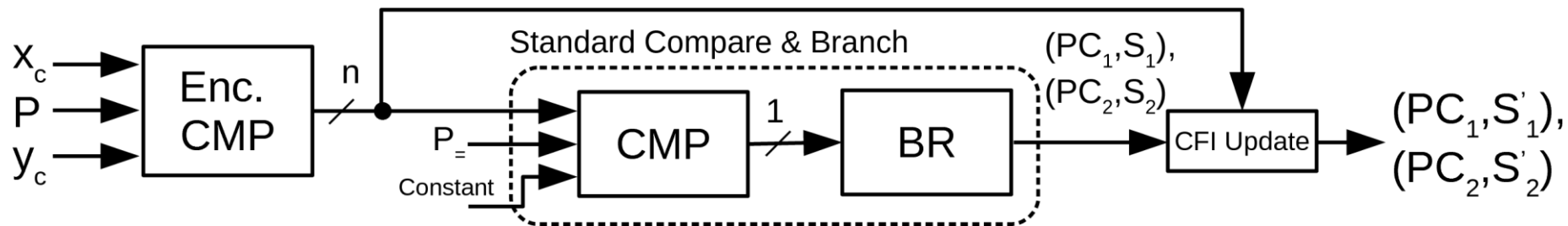
Step 2: Use a standard branch for the actual branching

Use a standard branch for the actual branching



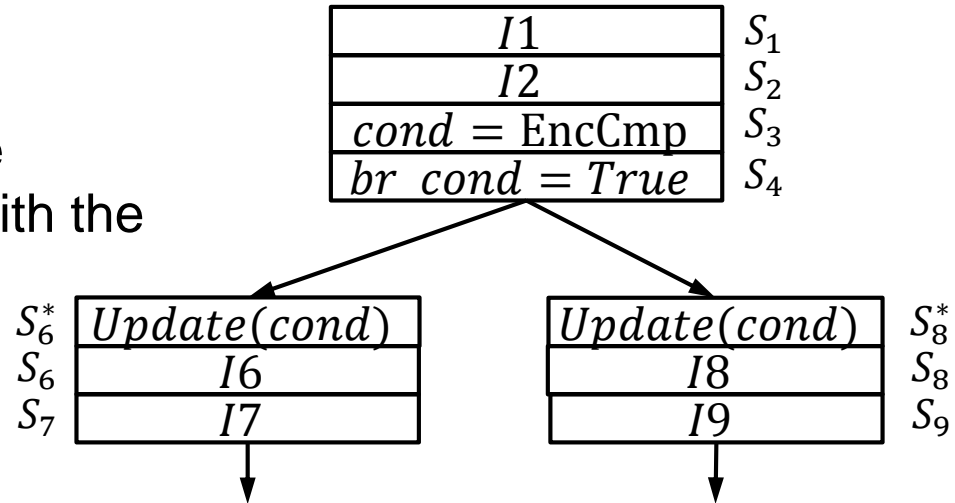
Step 3: Link Comparison Result to CFI State

... and link it to the CFI state.



Example: Protected Conditional Branch

1. Compute the encoded comparison
2. Perform a conditional branch
3. At the branch target: **Link** the redundant **condition** value with the CFI state



Wrong branch and wrong condition lead to invalid CFI state

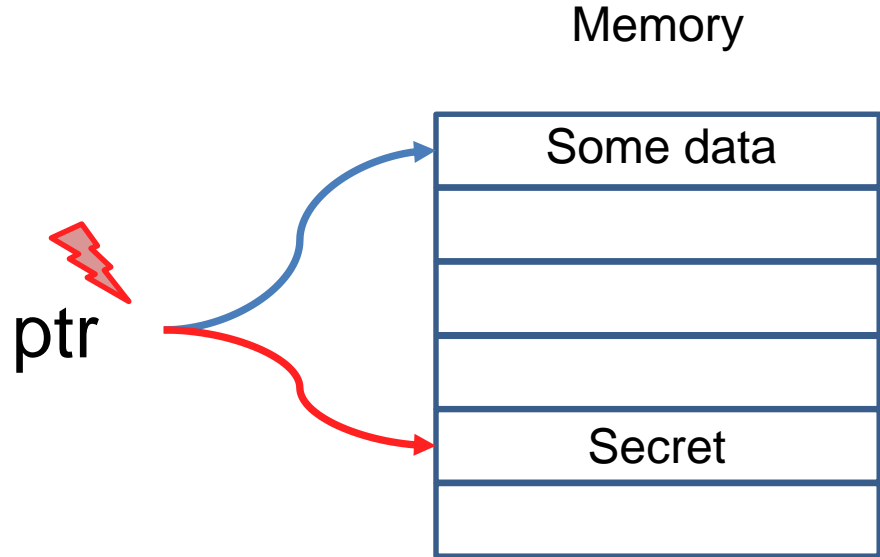
Prototype Evaluation

- Added new branch instruction to inject first operand to the CFI state
- LLVM-based toolchain
 - Automatically identifies conditional branches
 - Encodes dependent data-flow graph to AN-code domain
 - Inserts software-based comparison algorithm
- Overhead on par with state-of-the-art duplication approaches

Secure Memory Access

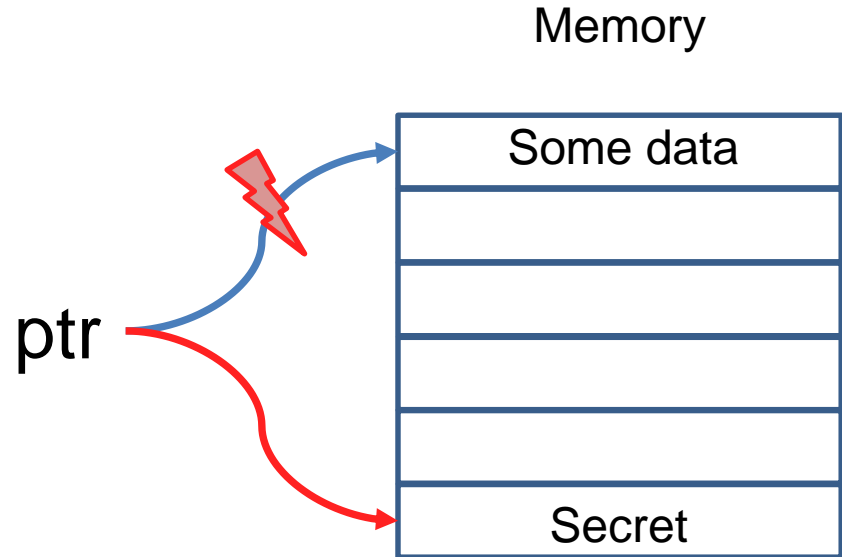
Faulting the Pointer

- Faulted pointer redirects the memory access



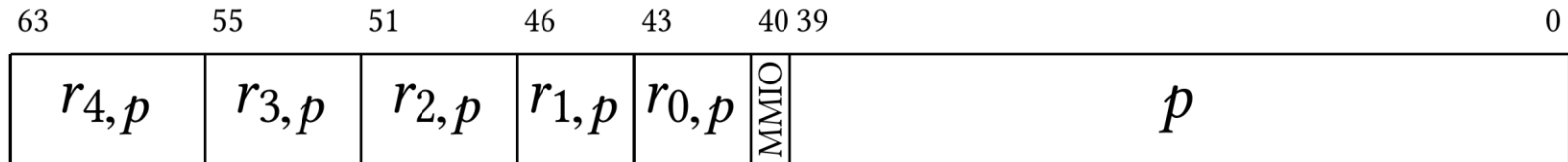
Faulting the Memory Access

- Faulted pointer redirects the memory access
- Faulting the memory access itself leads to a wrong access



Pointer Protection with Residue Codes

- Use multi-residue code to protect the pointer
 - Gives direct access to the functional value → no expensive decoding required
 - Supports pointer arithmetic
- Redundancy stored in the pointer

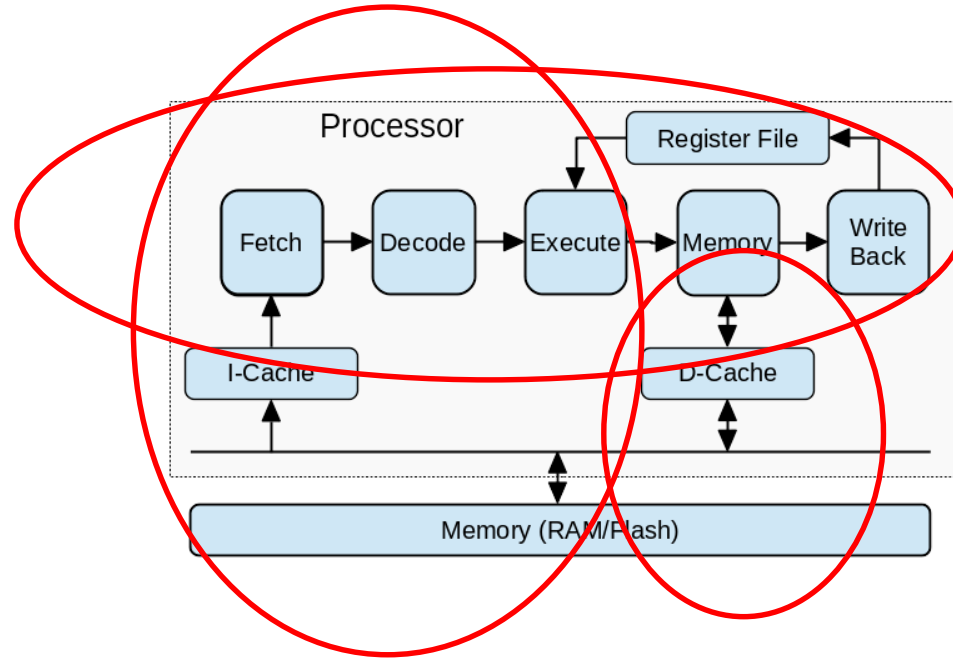


Protecting the Memory Accesses

- Write memory in the form $mem[p] = l_p(D_{Reg})$
- Inverse to read data back $D_{Reg} = l_p^{-1}(mem[p])$
- Xor operation \rightarrow chosen for low-overhead
 - $mem[p] = p \oplus D_{Reg}, \quad D_{Reg} = p \oplus mem[p]$
 - **Problems** with granularity
- Use a byte-wise linking granularity to support arbitrary accesses

Prototype Evaluation

- FPGA prototype based PULP by ETH Zurich with 5% area overhead
- ISA extension residue arithmetic and linked memory accesses
- Custom LLVM compiler prototype transforms all pointers
- Transformed all data pointers, protected all pointer arithmetic, replaced all memory accesses
- ~7% runtime and ~10% code size overhead





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